

A Novel Algorithm with Reduced Mutual Information for Smart Meter Privacy Protection

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Abstract—Fine-grained power utilization information from savvy meters might build the gamble of clients uncovering protection. This paper plans a technique to understand the stowing away of continuous power interest through family energy capacity gear (for example vehicle battery). By planning the rate contortion work issue and tackling it utilizing Blahut-Arimoto calculation, the accompanying beneficial impacts are accomplished: 1) the procedure is attainable; 2) the shared data between continuous power interest and preset meter information can be limited; 3) the information displayed in the meter are normal. Additionally, to confirm the viability of the proposed calculation, examinations with the conventional best-exertion and it are directed to step calculation. The outcomes show the proposed calculation has preferred execution over the over two calculations concerning the common data measurements. At last, the impact of most extreme battery power on the procedure is additionally mimicked.

Keywords—smart metering; privacy protection; load monitor; data analysis; rated distortion theory

I. INTRODUCTION

Due to its self-healing function, safety and compatibility, the smart grid has been developed by countries all over the world [1]-[3]. One of the foundations for efficient operation of smart grid is the Advanced Metering Infrastructure (AMI), which relies on smart meters. However, fine-grained electricity consumption data recorded by smart meter poses a risk of privacy leakage for users. For example, by using non-intrusive load monitoring data analysis techniques, users' usage patterns and personal habits can be obtained [4], [5]. Since the possibility of privacy leaks in smart meters, parts of the United States and Europe are prohibited from installation.

Many algorithms have been proposed for privacy-preserving of users' electricity consumption data recorded by smart meters [6]-[9]. Kalogridis et al. introduce a best-effort (BE) algorithm [10], which tries to keep the meter data fixed. Unfortunately, for the constraint of battery power and capacity, BE algorithm will definitely expose user's actual electricity data and thus leak privacy. Stephen proposes an non-intrusive load leveling (NIL) algorithm that controls the battery to charge/discharge when the capacity is too low/high to maintain a fixed value [11]. Given insights from the failure of the algorithms above to handle load peaks, a

different approach based on quantizing the demand load into a step function is proposed by [12].

However, none of these works have considered optimal control algorithms to protect the privacy of smart meter data. The optimal control strategy minimizes the impact of battery hardware constraints on the strategy while ensuring low mutual information.

Good protection can be achieved only if the data presented by the meter is preset and has a low correlation with the actual power demand. In this paper, an algorithm based on rated distortion theory that minimizes the mutual information between real-time power demand and the pre-set meter data is proposed.

The rest of the paper is organized as follows. In Section II, we introduce the system model. Section III defines the problem to be solved. A classical algorithm for solving rated distortion function is proposed in Section IV. Performance evaluation is proposed in Section V. The paper is concluded in Section VI.

II. SYSTEM MODEL

The considered input/output discrete time model is shown in Figure 1. Input X_t is the power demand of the home device at time t . Output Y_t is the energy obtained from the utility provider (UP). B_t is the power of the battery. The positive value indicates that the battery is charging, otherwise the battery is discharging. As described in the model, the real-time power demand X_t is jointly supplied by the battery power and the grid power, and EMU's (Energy Management Unit) adjustment of B_t can realize the change of Y_t . In the end, people can partially cover up the actual electricity consumption to protect privacy. The following part introduces the mathematical models of input and output power.

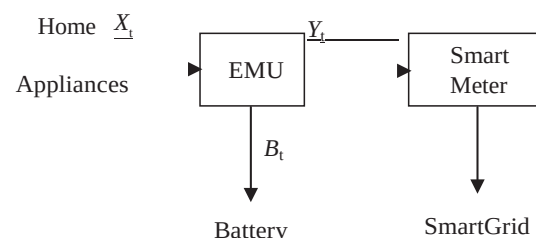


Figure 1. The system model

A. InputPowerModel

Definition1(HistorelectricitydatasetA):Thehistorical electricitydatasetof the family, denoted byA.

Definition2:(UniqueConsumptionSetG):Grepresentstheset ofunique consumptionrates inA.

Definition3(Functionh(.)):h(.)isdefinedtocalculate thenumberoftimesacertainpowerappearsinthefamilyhistorical electricity data.

Definition4(Functiong(.)):Thenumberofelements in asetcan becalculatedbyg(.).
Foreaseofanalysis,theinputpowerdemandisassumed

to beindependentandidenticallydistributed(i.i.d.)with

$$p(X=\omega) = \frac{h(\omega)}{g(A)}, \omega \in G.$$

B. OutputPowerModel

Toprotectprivacy,theelectricityconsumption Y_t recordedby hemetershouldbewhattheuserwantstopresent. Therefore, a data collection y is customized for Y_t to select. For the feasibility of customized set of meter data,someconstraintsneed tobeconsidered:

- Constraint1: $\max(y) \leq \max(A) + B_{\max}$
- Constraint2: $\min(y) \geq \min(A) - B_{\max}$
- Constraint3:Arrangeyinorder, andbareanytwo adjacent numbersiny, then: $a - b \leq 2B_{\max}$
- Constraint4: $g(y) \geq \frac{\max(A) - \min(A)}{2B_{\max}}$

If constraint 1 and 2 are not met, the maximum/minimumvalue of y can be directly deleted. Because it is impossible toachievebyadjusting thebattery power.

Constraint 3 and 4 are to ensure thatfor any X_t , at leastone element in y can be reached by charging or dischargingthebattery.

III. PROBLEMDEFINITION

In this section, an optimization framework is proposed tosearchtheoptimalstrategythatminimizesthemutual informationbetweeninputpowerdemandandpresetmeterdata. Ratedistortiontheoryhasbeenproventobevery effectiveincalculatingtheminimummutualinformationbetween woprobabilitydistributionsetsundergiven constraints.Tothisend,aratedistortionfunctionproblemisformulatedtocapturethebestfeasiblestrategy.

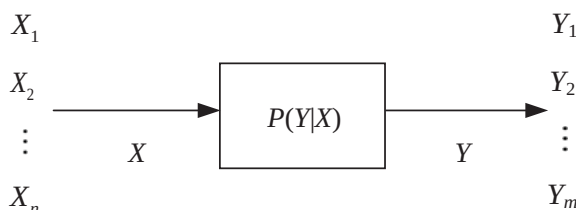


Figure2.Informationtransmissionmodel

A. SystemDynamics

Since the uncertainty of real-time electricity consumptionmakes the quantization problem of solving a single randomvariableverycomplicated,a jointdescriptionismoreeffective than a single description. Based on rate distortiontheory, an information transmission model is developed todescribepowerdemandandmeterdatajointly,whichisshownin Figure2.XinFigure2standsforreal-timepower demandandYisthecustomizedmeterdataset. $p(Y|X)$ in thefigurerepresentshetransmissionprobabilityofthechannel, whichisrepresented bya matrixas:

$$[p(Y|X)] = \begin{bmatrix} p(Y_1|X_1) & p(Y_2|X_1) & \cdots & p(Y_m|X_1) \\ p(Y_1|X_2) & p(Y_2|X_2) & \cdots & p(Y_m|X_2) \\ \vdots & \vdots & \ddots & \vdots \\ p(Y_1|X_n) & p(Y_2|X_n) & \cdots & p(Y_m|X_n) \end{bmatrix}$$

Each element of the matrix represents the probability thatthe real-time power demand is X_i and the meter data is Y_i .After setting the meter data, all strategies can be presented intheformofthismatrix.

B. DistortionFunction

After fully considering the role of battery power in thestrategy, absolute distortion is used as the distortion functionofthispaper.

$$d(X, Y) = \begin{cases} |X_t - Y_t|, & |X_t - Y_t| \leq B_{\max} \\ B_{\max}, & \text{otherwise} \end{cases} \quad (1)$$

where $B_{\max}$ is the maximum distortion. For the data that cannot be achieved by charging or discharging the battery, the distortion should be large enough so that there will be no unreachable meter data be selectedgiven the expected distortion.

C. ObjectiveFunction

Theorem1[13].Forthesourceofi.i.d.,ifthedistributionis $p(x)$ and thebounded distortionfunctionis $d(x, \hat{x})$, then the ratedistortionfunctionisequal tothe correspondinginformationratedistortionfunction,whichis:

$$R(D) = R^{(I)}(D) = \min_{p(\hat{x}|x): \sum_{(x, \hat{x})} p(x)p(\hat{x}|x)d(x, \hat{x}) \leq D} I(X, \hat{X}) \quad (2)$$

Accordingtotheorem1,thefollowingobjectivefunction canbeestablished:

$$\begin{aligned} R(D) &= \min_{p(Y|X): \sum_{(X, Y)} p(X)p(Y|X)d(X, Y) \leq D} I(X, Y) \\ \text{s.t. } &\begin{cases} \sum_i p(Y|X) = 1 \\ p(Y|X) \geq 0 \end{cases} \end{aligned} \quad (3)$$

Theintermediatevariable D can be shown as:

$$\begin{aligned} D &= \sum_{i=1}^{nm} \sum_{j=1}^{nm} p(X_i Y_j) d(X_i, Y_j) \\ &= \sum_{i=1}^n \sum_{j=1}^m p(X_i) p(Y_j|X_i) d(X_i, Y_j) \end{aligned} \quad (4)$$

Disnotonlytheexpectederror,butalsoequalsinvaluetothetheexp
ectedpowerofthebattery.Theminimumvalueis
takenfromalltransitionprobabilitymatrices $p(Y|X)$ that
makethe joint distribution $p(X)p(Y|X)$ satisfy the
expectationofdistortionconstraint D .Forthisproblem,various
toolsof rate distortiontheorycanbe utilized.

IV. SOLUTIONMETHODOLOGY

Thissectionintroducesaclassicalmethod,Blahut-
Arimotoalgorithm[13],forcalculatingtheratedistortionfunction.
A description of the Blahut-Arimoto algorithm is presented in
Algorithm 1. It is executed by the EMU to
learntheoptimalstrategybasedonhistoricalhouseholdconsumpti
on data.

TheBlahut-Arimotoalgorithm uses an alternating
minimizationalgorithmto calculatetheminimummutual
informationbetweenpowerdemand(X)andcustomizedmeterdata
(Y).Ineachiteration,foragiveninput
distribution $p(X)$ and output distribution $p(Y)$, first
calculatethe $p(Y|X)$ thatminimizesmutualinformationunder
distortionconstraints:

$$p(Y|X) = \frac{p(Y)e^{-\lambda d(X,Y)}}{\sum p(Y)e^{-\lambda d(X,Y)}} \quad (5)$$

where λ

isthelagrangeparameter.

Lemma1[13]:Let $p(X)p(Y|X)$ bethegivenjointdistribution.
Then the distribution $r^*(Y)$ which minimizes
therelativeentropy

$$D(p(X)p(Y|X)||p(X)r(Y)) \\ = \min_{r(Y)} (D(p(X)p(Y|X)||p(X)r(Y))) \quad (6)$$

isthemarginaldistribution $r^*(Y)$ correspondingtop $p(Y|X)$,

where $r^*(Y) = \sum p(X)p(Y|X)$.

Accordingtothelemma1,theoutputdistribution $p(Y)$,
whichminimizesmutualinformation,canbecalculatedasfollows:

$$p(Y) = \sum_x p(X)p(Y|X) \quad (7)$$

Theentropyandjointentropyoftherandomvariableinstep 8
canbe obtainedbythefollowingtwoformulas:

$$H(X) = -\sum p(X_i) \log_2 p(X_i) \quad (8)$$

$$H(X,Y) = -\sum \sum p(X_i, Y_j) \log_2 p(X_i, Y_j) \quad (9)$$

$I_1(X;Y)$ and $I_2(X;Y)$ inStep10representthemutual

informationobtainedbytwoadjacentcomputations.

As the number of iterations increases, this
minimizationprocessconvergesto $R(D)$.

V. PERFORMANCEEVALUATION

In this section, the proposed Battery-based Load
Hiding(BLH) strategy is compared with two other common
BLH algorithms, BE andStepping, to evaluate its
performance.Moreover,theinfluenceofthemaximumpowerofth
ebatteryonthisstrategyisalsostudiedandshownintheFigure4.

Algorithm1.Blahut-ArimotoAlgorithm

Step1:Inputthedistributionofelectricitydemand $p(X)$.
Step2:Inputiterationprecision $prec$,maximumnumber
ofiterationstandLagrangeparameter λ .

Step3:Inputthepresetmeterdataanditsdistribution
 $p(Y)$.

Step4: Repeat(for each λ):

Step5: $c=1$

Step6:Calculatethetransitionprobabilitymatrix $p(Y|X)$ from (5).

Step7:Obtaintheoutput

distribution $p(Y)$ from(7).Step8:Calculatemutualinformati
onbetween X,Y :

$$I(X;Y) = H(X) + H(Y) - H(X,Y)$$

Step9: $c = c+1$

Step10:Until $c > t$ or $I_1(X;Y) - I_2(X;Y) \leq prec$

Step11:Output $p(Y|X)$

BEalgorithm

BEalgorithmtriestoresist(tothedegreepossible)

againstpowerloadchanges.Whenthe requiredloadiseitherlargeror
smallerrespectivelythanthe previouslymetered
load,EMUwillchargeordischarge the battery tomakeup the differen
ce.

Steppingalgorithm

Steppingalgorithmaimstoquantifythedemandloadinto
astepfunction,whichisdeterminedbythebattery's
maximumcharge/discharge rate.In thesteppingalgorithm,

theEMUforcestheexternalloadtobemultiplesof $\mathcal{E}$,which
isdeterminedbythebattery's parameters.

A. InfluenceofDifferentAlgorithms

Weconsidera24-hourplanningperiodinthesimulation.
Thesmartmetercurrentlydeployedinprivatedwellingsin
theEastMidlands,UK.Tounderstandtheeffectsofthesealgorithmsi
ntuitively,theoriginalelectricitydataisalso
shown.Themaximumpowerofthebatteryusedforthe
simulationis1000Wandcapacityis0.5kWh.Themeterdata

presetbytheproposedalgorithmis(0,200,400,800,1000,
1200,1600,2000,2400,2800,3200,3600, 4000,4800).

The mutual information between the current meter data and the
al-time power demand is presented by using the
upperrightcorneroftheexternalloadfigureforeachalgorithm.

SimulationresultsshowthattheBEalgorithmhasagood
effect on the power demand less than the maximum power
ofthe battery. But, due to the constraint of maximum power
ofthe battery, it has limited ability to hide the large real-
timeelectricity consumptionrate. Itcanalso
beseenfromthefigure that since the stepping algorithm
maximizes the errorbetween the power demand and the meter
data, the battery isvery easy to overcharge or over discharge.
Eventually lead totheleakageofreal-
timeelectricityinformation.Ouralgorithmnotonlyhasagoodeffe
ctonthe power demandlessthanthemaximumpowerofthebattery,
butalsocanfind a good representation of the power demand
greater thanthe maximumpowerofthebatteryfromthe preset
powerset.

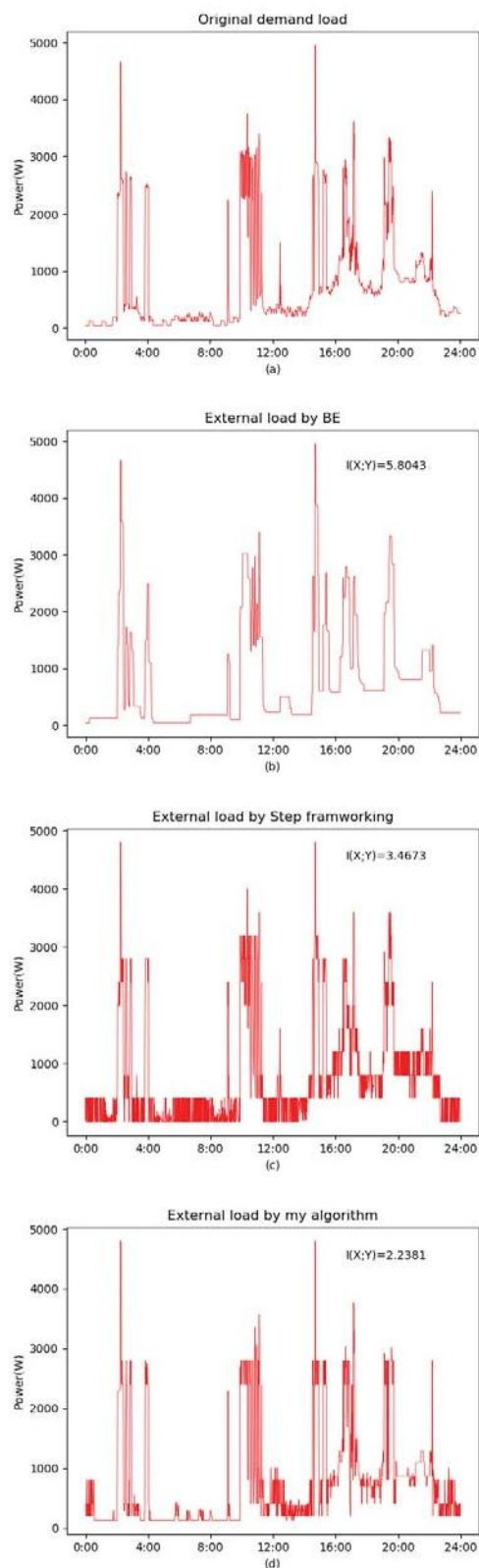


Figure3.Influenceof differentialalgorithms

B. Influenceof Different Batteries

Figure 4 shows the impact of different maximum power batteries on this strategy. The abscissa is the expected error (which is numerically equal to the expected power of the battery) and the ordinate is the minimum mutual information. It is a non-increase convex function on D , indicating that the minimum mutual information can be achieved by getting smaller and smaller as the expected power of the battery increases.

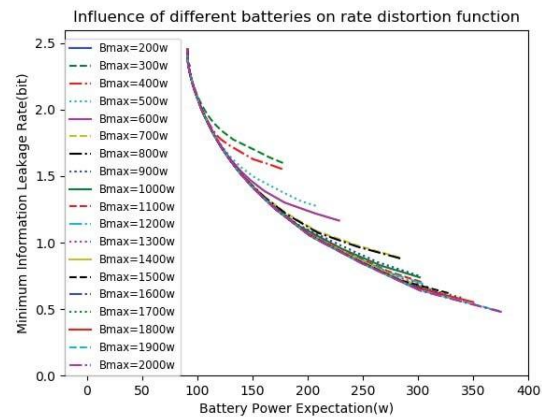


Figure4.Influenceof different batteries on rated distortion function

VI. CONCLUSION

The existing BLH strategies do not consider optimal control algorithms to protect smart meter data. We present the data presented by the meter and propose a new BLH algorithm with reduced mutual information based on the theory of rate distortion. Compared with other strategies, the algorithm proposed in this paper is an online optimization algorithm designed for the family's historical power consumption situation and user needs, so it can achieve better protection effects. The transition probability matrix is calculated in advance, so the smart meter does not need to carry out a lot of calculations in real time, and the hardware requirements of the smart meter are not high. The mutual information is used as a measure to compare the proposed algorithm with BE and Stepping algorithm, and the experimental results demonstrate the effectiveness of the proposed algorithm. In the future, the impact of other factors, such as battery capacity, on this strategy will be studied.

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